50 Ways to Leak Your Data: An Exploration of Apps' Circumvention of the Android Permissions System

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is crucial to protect this information from unauthorized ac-

Abstract

Modern smartphone platforms implement permission-based models to protect access to sensitive data and system resources. However, apps can circumvent the permission model and gain access to protected data without user consent by us ing both covert and side channels. Side channels present in the implementation of the permission system allow apps to access protected data and system resources without permission; whereas covert channels enable communication between two colluding apps so that one app can share its permissionprotected data with another app lacking those permissions. Both pose threats to user privacy.

In this work, we make use of our infrastructure that runs hundreds of thousands of apps in an instrumented environment. This testing environment includes mechanisms to monitor apps' runtime behaviour and network traf c. We look for evidence of side and covert channels being used in practice by searching for sensitive data being sent over the network for which the sending app did not have permissions to access in their software for things like crash reporting, development it. We then reverse engineer the apps and third-party libraries support, analytics services, social-network integration, and adresponsible for this behaviour to determine how the unauthovertising [16,62]. By design, any third-party service bundled methods to measure the static prevalence of the technique that we discover among other apps in our corpus.

Using this testing environment and method, we uncovered a embedded in that app can as well. number of side and covert channels in active use by hundreds In practice, security mechanisms can often be circumof popular apps and third-party SDKs to obtain unauthorized vented;side channelandcovert channelare two common access to both unique identi ers as well as geolocation datatechniques to circumvent a security mechanism. These chan-We have responsibly disclosed our ndings to Google and nels occur when there is an alternate means to access the prohave received a bug bounty for our work.

1 Introduction

tected resource that is not audited by the security mechanism, thus leaving the resource unprotected ide channel exposes a path to a resource that is outside the security mechanism; this can be because of a aw in the design of the security mechanism or a aw in the implementation of the design. A

Smartphones are used as general-purpose computers anothessic example of a side channel is that power usage of hardtherefore have access to a great deal of sensitive system novare when performing cryptographic operations can leak the sources (e.g., sensors such as the camera, microphone, **p**articulars of a secret key [42 GPS), private data from the end user (e.g., user email or con-

tacts list), and various persistent identi ers (e.g., IMEI). It

A covert channels a more deliberate and intentional effort We studied more than 88,000 apps across each category between two cooperating entities so that one with access torom the U.S. Google Play Store. We found a number of side some data provides it to the other entity without access toand covert channels in active use, responsibly disclosed our the data in violation of the security mechanism [43]. As an ndings to Google and the U.S. Federal Trade Commission example, someone could execute an algorithm that alternate(FTC), and received a bug bounty for our efforts. between high and low CPU load to pass a binary message to In summary, the contributions of this work include: another party observing the CPU load.

The research community has previously explored the potential for covert channels in Android using local sockets and shared storage [49], as well as other unorthodox means, such as vibrations and accelerometer data to send and receive data between two coordinated apps [3]. Examples of side channels include using device sensors to infer the gender of the user [51]or uniquely identify the user [72]. More recently, researchers demonstrated a new permission-less device ngerprinting technique that allows tracking Android and iOS devices across the Internet by using factory-set sensor calibration details [90]. However, there has been little research in detecting and measuring at scale the prevalence of covert and side channels in apps that are available in the Google Play Store. Only isolated instances of malicious apps or libraries inferring users' locations from WiFi access points were reported, a side channel that was abused in practice and resulted in about a million dollar ne by regulators [82].

In fact, most of the existing literature is focused on under standing personal data collection using the system-supported access control mechanisms (i.e., Android permissions). With increased regulatory attention to data privacy and issues sur rounding user consent, we believe it is imperative to under stand the effectiveness (and limitations) of the permission system and whether it is being circumvented as a preliminary step towards implementing effective defenses.

To this end, we extend the state of the art by developing methods to detect actual circumvention of the Android per mission system, at scale in real apps by using a combination of dynamic and static analysis. We automatically executed over 88,000 Android apps in a heavily instrumented environment with capabilities to monitor apps' behaviours at the system and network level, including a TLS man-in-the-middle proxy. In short, we ran apps to see when permission-protected data was transmitted by the device, and scanned the apps to see which oneshould nothave been able to access the transmitted data due to a lack of granted permissions. We grouped our ndings by whereon the Internet/what data type was sent, as this allows us to attribute the observations to the actual app developer or embedded third-party libraries. We then reverse engineered the responsible component to determine exactly how the data was accessed. Finally, we statically analyzed our entire dataset to measure the prevalence of the channel. We focus on a subset of thengerouspermissions that prevent apps from accessing location data and identi ers. Instead of imagining new channels, our work focuses on tracing evidence that suggests that side- and covert-channel abuse is occurring in practice.

We designed a pipeline for automatically discovering vulnerabilities in the Android permissions system through a combination of dynamic and static analysis, in effect creating a scalable honeypot environment. We tested our pipeline on more than 88,000 apps and discovered a number of vulnerabilities, which we responstudy, including the side and covert channels we discovered he kernel directly as well. For some permission-protected and their prevalence in practice. Section describes related resources, such as network sockets, the reference monitor is work. Section 6 discusses their potential legal implications. the kernel, and the request for such resources bypasses the Section7 discusses limitations to our approach and concludesplatform framework and directly contacts the kernel. Our with future work. work discusses how real-world apps circumvent these system

2 Background checks placed in the kernel and the platform layers. The Android permissions system serves an important pur

pose: to protect users' privacy and sensitive system resources The Android permissions system has evolved over the years from deceptive, malicious, and abusive actors. At the very from an ask-on-install approach to an ask-on- rst-use approach. While this change impacts when permissions are not be able to access data protected by that permissions are

granted and how users can use contextual information to real practice, this is not always the case. son about the appropriateness of a permission request, the backend enforcement mechanisms have remained largely un2.2 Circumvention

changed. We look at how the design and implementation of

the permission model has been exploited by apps to bypass ent ways [317,49,51,52,54,70,72,74]. The use of covert and these protections.

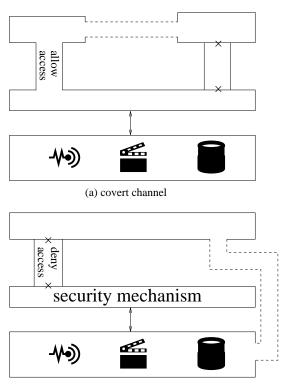
Android Permissions 2.1

Android's permissions system is based on the security prin-Commission (FTC) has ned mobile developers and thirdciple of least privilege. That is, an entity should only have the minimum capabilities it needs to perform its task. This party libraries for exploiting side channels: using the MAC adstandard design principle for security implies that if an app acts maliciously, the damage will be limited. Developers must declare the permissions that their apps need beforehand, and the user is given an opportunity to review them and decide whether to install the app. The Android platform, however, does not judge whether the set of requested permissions are all strictly necessary for the app to function. Developers are free to request more permissions than they actually need and users are expected to judge if they are reasonable.

The Android permission model has two important aspects: obtaining user consent before an app is able to access any of its requested permission-protected resources, and then ensur ing that the app cannot access resources for which the user has not granted consent. There is a long line of work uncovering issues on how the permission model interacts with the user: users are inadequately informed about why apps need permissions at installation time, users misunderstand exactly what the purpose of different permissions are, and users lack context and transparency into how apps will ultimately use their granted permissions [230,78,86]. While all of these are critical issues that need attention, the focus of our work is to understand how apps are circumventing system checks to verify that apps have been granted various permissions.

When an app requests a permission-protected resource, the resource manager (e.d. cati on Manager, Wi Fi Manager, etc.) contacts thectivityServiceManager, which is the reference monitoin Android. The resource request originates from the sandboxed app, and the nal veri cation happens inside the Android platform code. The platform is a Java oper ating system that runs in system space and acts as an interface for a customized Linux kernel, though apps can interact with

side channels, however, is particularly troublesome as their usage indicates deceptive practices that might mislead even diligent users, while underscoring a security vulnerability in



(b) side channel

Figure 1: Covert and side channels. (a) A security mechanism allowsapp1 access to resources but denips2 access; this is circumvented by app2 using app1 as a facade to obtain access over a communication channel not monitored by the security mechanism. (b) A security mechanism denips1 access to resources; this is circumvented by accessing the resources through a side channel that bypasses the security mechanism.

being protected by the same permission. A classical example of a side channel attack is the timing attack to ex Itrate an encryption key from secure storage [42]. The system under attack is an algorithm that performs computation with the key and unintentionally leaks timing information—i.e., how long it runs—that reveals critical information about the key.

Side channels are typically an unintentional consequence of a complicated system. ("Backdoors" are intentionally-created side channels that are meant to be obscure.) In Android, a large and complicated API results in the same data appearing in different locations, each governed by different access control mechanisms. When one API is protected with permissions, another unprotected method may be used to obtain the same data or an ersatz version of it.

2.3 App Analysis Methods

Researchers use two primary techniques to analyze app behaviour: static and dynamic analysis. In short, static analysis studiessoftware as datay reading it; dynamic analysis studiessoftware as codby running it. Both approaches have the

goal of understanding the software's ultimate behaviour, but they offer insights with different certainty and granularity: static analysis reports instances of hypothetical behaviour; dynamic analysis gives reports of observed behaviour.

Static Analysis Static analysis involves scanning the code

requires building an instrumentation framework for possible behaviours of interest priori and then engineering a system to manage the endeavor.

Nevertheless, some apps are resistant to being audited when run in virtual or privileged environments [128]. This has led to new auditing techniques that involve app execution on real phones, such as by forwarding traf c through a VPN in order to inspect network communications [44,63]. The limitations of this approach are the use of techniques robust to man-in-the-middle attacks [281,61] and scalability due to the need to actually run apps with user input.

A tool to automatically execute apps on the Android platform is the Ul/Application Exerciser Monkey [6]. The Monkey is a UI fuzzer that generates synthetic user input, ensuring that some interaction occurs with the app being automatically tested. The Monkey has no context for its actions with the UI, however, so some important code paths may not be executed due to the random nature of its interactions with the app. As a result, this gives a lower bound for possible app behaviours, but unlike static analysis, it does not yield false positives.

Hybrid Analysis Static and dynamic analysis methods complement each other. In fact, some types of analysis benet from a hybrid approach, in which combining both methods can increase the coverage, scalability, or visibility of the analyses. This is the case for malicious or deceptive apps that actively try to defeat one individual method (e.g., by using obfuscation or techniques to detect virtualized environments or TLS interception). One approach would be to rst carry out dynamic analysis to triage potential suspicious cases, based on collected observations, to be later examined thoroughly using static analysis. Another approach is to rst carry out static analysis to identify interesting code branches that can then be instrumented for dynamic analysis to con rm the ndings.

3 Testing Environment and Analysis Pipeline

Our instrumentation and processing pipeline, depicted and described in Figure

3.1 App Collection

We wrote a Google Play Store scraper to download the mostpopular apps under each category. Because the popularity distribution of apps is long tailed, our analysis of the 88,113 most-popular apps is likely to cover most of the apps that people currently use. This includes 1,505 non-free apps we pur chased for another study [38]. We instrumented the scraper to inspect the Google Play Store to obtain application executables (APK les) and their associated metadata (e.g., number of installs, category, developer information, etc.).

As developers tend to update their Android software to add

After running the app, the kernel, platform, and network personal information encoded in network ows, such as gzip, logs are collected. The app is then uninstalled along with anybase64, and ASCII-encoded hexadecimal. Additionally, we other app that may have been installed through the process of earch for personal information directly, as well as the MD5, automatic exploration. We do this with a white list of allowed SHA1, and SHA256 hashes of it.

apps; all other apps are uninstalled. The logs are then cleared After analyzing thousands of network traces, we still nd and the device is ready to be used for the next test.

3.3 Personal Information in Network Flows

Detecting whether an app has legitimately accessed a given reof reverse engineering and static analysis to understand the source is straightforward: we compare its runtime behaviourprecise technique. We frequently found a further use of AES with the permissions it had requested. Both users and reencryption applied to the payload before sending it over the searchers assess apps' privacy risks by examining their renetwork, often with hard-coded AES keys.

quested permissions. This presents an incomplete picture, A few libraries followed best practices by generating ranhowever, because it only indicates what data amaight acdom AES session keys to encrypt the data and then encrypt cess, and says nothing about with whom it may share it and he session key with a hard-coded RSA public key, sending under what circumstances. The only way of answering these both the encrypted data and encrypted session key together. questions is by inspecting the apps' network traf c. However, To de-cipher their network transmissions, we instrumented identifying personal information inside network transmissions the relevant Java libraries. We found two examples of thirdrequires signi cant effort because apps and embedded thirdparty SDKs "encrypting" their data by XOR-ing a keyword party SDKs often use different encodings and obfuscationover the data in a Viginère-style cipher. In one case, this was techniques to transmit data. Thus, it is a signi cant technical in addition to both using standard encryption for the data challenge to be able to de-obfuscate all network traf c and using TLS in transmission. Other interesting approaches insearch it for personal information. This subsection discussescluded reversing the string after encoding it in base64 (which how we tackle these challenges in detail. we refer to as "46esab"), using base64 multiple times (base-

Personal Information We de ne "personal information" as any piece of data that could potentially identify a speci c ings and our entire dataset is then re-analyzed. individual and distinguish them from another. Online compa-

nies, such as mobile app developers and third-party advertis 3.4 Finding Side and Covert Channels ing networks, want this type of information in order to track

users across devices, websites, and apps, as this allows the pance we have examples of transmissions that suggest the to gather more insights about individual consumers and thuspermission system was violated (i.e., data transmitted by an generate more revenue via targeted advertisements. For thispp that had not been granted the requisite permissions to reason, we are primarily interested in examining apps' access to so), we then reverse engineer the app to determine how it to the persistent identi ers that enable long-term tracking, as circumvented the permissions system. Finally, we use static well as their geolocation information. analysis to measure how prevalent this practice is among the

We focus our study on detecting apps using covert and siderest of our corpus. channels to access speci c types of highly sensitive data, in-

cluding persistent identi ers and geolocation information. No- Reverse Engineering After nding a set of apps exhibittably, the unauthorized collection of geolocation information ing behaviour consistent with the existence of side and covert in Android has been the subject of prior regulatory action [82]. channels, we manually reverse engineered them. While the Table1 shows the different types of personal information that reverse engineering process is time consuming and not easily we look for in network transmissions, what each can be used automated, it is necessary to determine how the app actually for, the Android permission that protects it, and the subsecobtained information outside of the permission system. Betion in this paper where we discuss ndings that concern side cause many of the transmissions are caused by the same SDK and covert channels for accessing that type of data. code, we only needed to reverse engineer each unique

Decoding Obfuscations In our previouswork [66], we found instances of apps and third-party libraries (SDKs) using obfuscation techniques to transmit personal information over the network with varying degrees of sophistication. To identify and report such cases, we automated the decoding of a standard suite of standard HTTP encodings to identify

(sa4b6e). Each new discovery is added to our suite of decod-

base6464), and using a permuted-alphabet version of base64

new techniques SDKs and apps use to obfuscate and encrypt network transmissions. While we acknowledge their effort to protect users' data, the same techniques could be used to hide deceptive practices. In such cases, we use a combination

Table 1: The types of personal information that we search for, the permissions protecting access to them, and the purpose for which they are generally collected. We also report the subsection in this paper where we report side and covert channels for accessing each type of data, if found, and the number of apps exploiting each. The dynamic column depicts the number of apps that we directly observed inappropriately accessing personal information, whereas the static column depicts the number of apps containing code that exploits the vulnerability (though we did not observe being executed during test runs).

Data Type	Permission	Purpose/Use	Subsection	N ^o of Apps		N ^o of SDKs		Channel Type	
				Dynamic	Static	Dynamic	Static	Covert	Side
IMEI	READ_PHONE_STATE	Persistent ID	4.1	13	159	2	2	2	0
Device MAC	ACCESS_NETWORK_STATE	Persistent ID	4.2	42	12,408	1	1	0	1
Email	GET_ACCOUNTS	Persistent ID	Not Found						
Phone Number	READ_PHONE_STATE	Persistent ID	Not Found						
SIM ID	READ_PHONE_STATE	Persistent ID	Not Found						
Router MAC	ACCESS_WIFI_STATE	Location Data	4.3	5	355	2	10	0	2
Router SSID	ACCESS_WIFI_STATE	Location Data	Not Found						
GPS	ACCESS_FINE_LOCATION	Location Data	4.4	1	1	0	0	0	1

which data sources. For some particular apps and libraries, our work also necessitated reverse engineering C++ code; we used IdaPro [1] for that purpose.

The typical process was to search the code for strings cor responding to destinations for the network transmissions and other aspects of the packets. This revealed where the data was already in memory, and then static analysis of the code revealed where that value rst gets populated. As intentionallyobfuscated code is more complicated to reverse engineer, we Android protects access to the phone's IMEI with the $\ensuremath{\mathsf{READ}}\xspace_$

4.2 Network MAC Addresses

The Media Access Control Address (MAC address) is a 6-byte identi er that is uniquely assigned to the Network Interface Controller (NIC) for establishing link-layer communications. However, the MAC address is also useful to advertisers and analytics companies as a hardware-based persistent identi er, similar to the IMEI.

Android protects access to the device's MAC address with the ACCESS_NETWORK_STATE permission. Despite this, we observed apps transmitting the device's MAC address without

Table 2: SDKs seen sending router MAC addresses and also containing code to access the ARP cache. For reference, we report the number of apps and a lower bound of the total number of installations of those apps. We do this for all apps containing the SDK; those apps that nothaveAccess_wiFi_STATE, which means that the side channel circumvents the permissions system; and those apps which do have a location permission, which means that the side channel circumvents location revocation.

SDK Name	Contact	Incorporation	Total	Prevalance	Wi-Fi	Permission	No Loca	ation Permission
	Domain	Country	(Apps)	(Installs)	(Apps)	(Installs)	(Apps)	(Installs)
AlHelp Huq Industries	cs30.net huq.io	United States huq.io	30	334 million	3	210 million	12	195 million

from the photo library, which included the phone's precise location in its exchangeable image le format (EXIF) data. The app actually processed the image le: it parsed the EXIF metadata—including location—into a JSON object with labelled! ati tude and! ongi tude elds and transmitted it to their server.

While this app may not be intending to circumvent the permission system, this technique can be exploited by a malicious actor to gain access to the user's location. Whenever a new picture is taken by the user with geolocation enabled, any app with read access to the photo library (i.e., READ_EXTERNAL_STORAGE) can learn the user's precise location when said picture was taken. Furthermore, it also allows obtaining historical geolocation xes with timestamps from the user, which could later be used to infer sensitive information about that user.

5 Related Work

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other permissions that, while not labeleddasgerous, can still give access to sensitive user data. One example is the BLUETOOTH

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